

## INTRODUCTION

What exactly is a quantum random walk?

- The term was first coined in the 1993 paper by Aharonov, Davidovich and Zaguray[1], although Feynman[2] had introduced the idea of a quantum walk many years before in 1940, except he did not call it a quantum walk
- There is both a discrete and continuous quantum random walk, just as in the classical case
- They can also be applied to an undirected graph  $G(V, E)$ , which means we can implement them as quantum algorithms
- The behaviour of both the discrete and continuous quantum random walks is very different from the classical case
- They do not have gaussian properties and they do not converge to limiting distributions, which means it is very hard to analytically calculate their statistical properties
- The width(standard deviation) of both the distributions on the line is found to be  $\sigma_{Quant} = t$ , which is much greater than the classical case of  $\sigma_{Class} = \sqrt{t}$
- Which means the walker travels much further on average from their starting position!
- It is this property of the walks that has sparked a lot of excitement, especially in regards to algorithmic development for quantum computers

## DISCRETE QUANTUM WALK

The Set Up

- To implement the walk on the line we introduce a coin space  $\mathcal{H}_c$  and a position space  $\mathcal{H}_p$ , so our total space is  $\mathcal{H} = \mathcal{H}_p \otimes \mathcal{H}_c$ , where we must introduce a unitary operator  $\hat{U}$  that acts on  $\mathcal{H}$
- Where the coin space represents the states of a coin and is spanned by the basis vectors  $\{|R\rangle, |L\rangle\}$ . The position space is spanned by the position states  $\{|x\rangle : x \in \mathbb{Z}^{+,0}\}$
- We first prepare the system in some initial state  $|\psi(0)\rangle = |0\rangle \otimes |\phi(0)\rangle$ , where  $|\phi(0)\rangle = |L\rangle, |R\rangle$  or a superposition of the two
- We next introduce a translational shift operator  $\hat{S} : |x+1\rangle \otimes |R\rangle$ , which only acts on  $\mathcal{H}_p$
- We then introduce our coin operator  $\hat{C}$ . If we would like the symmetric quantum walk on the line, we can choose  $\hat{C} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & i \\ i & 1 \end{pmatrix}$
- This can also be done via a different coin  $\hat{C}$  and changing our initial condition  $|\phi(0)\rangle$  to generate the same effect.

The Walk

- So the dynamics of the walk is described by the unitary acting on the initial state  $|\psi(0)\rangle$

$$|\psi(t+1)\rangle = \hat{U}^t |\psi(t)\rangle = (\hat{S}(\hat{C} \otimes \hat{I}))^t |\psi(0)\rangle \quad (1)$$

- which leads to the symmetrical, non-Gaussian distribution in figure 2

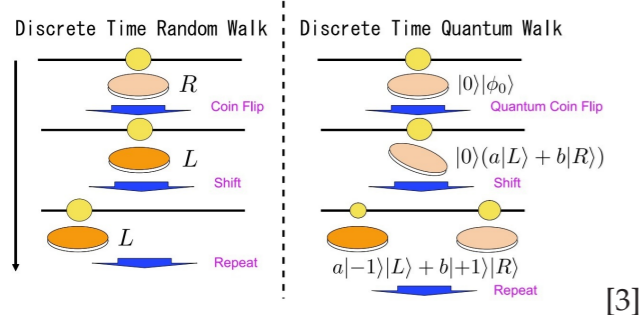


Figure 1: A pictorial representation of how the discrete quantum walk evolves. As more iterations are carried out, the “walker” can be in a superposition of multiple position states. This is ultimately why the quantum walker travels so much further.

## CONTINUOUS RANDOM WALK

Set Up

- The continuous time random walk does not have the extra structure of the coin space  $\mathcal{H}_c$ , instead its Hilbert space  $\mathcal{H}$  is simply the position space  $\mathcal{H}_p$
- As in the continuous Markov process, the time is also continuous and takes place on the graph  $G(V, E)$ , except the vertices are now the quantum states  $\{|1\rangle, \dots, |v\rangle\}$
- We aim to turn our transition matrix  $M$ , in to a unitary operator

$$M_{ab} = \begin{cases} k\gamma, & a = b \text{ and } k \text{ is the degree of vertex } a \\ -\gamma, & a \neq b, a \text{ and } b \text{ are connected by an edge} \\ 0, & a \neq b, a \text{ and } b \text{ are not connected} \end{cases} \quad (2)$$

The Walk

- We do this by multiplying the transition matrix by the imaginary unit  $i$  [4], which then turns our transition matrix  $\hat{M}$  into the Hamiltonian  $\langle a|\hat{H}|b\rangle = M_{ab}$
- Solving the Schrödinger equation

$$i \frac{d}{dt} \langle a|\psi(t)\rangle = \sum_b \langle a|\hat{H}|b\rangle \langle b|a\rangle \quad (3)$$

for a given Hamiltonian, with initial condition  $M_{ab} = \delta_{ab}$

- We gain the following unitary which describes the dynamics of the walk,  $\hat{U} = e^{-i\hat{H}t}$
- Thus our walk evolves, for some initial state  $|\psi(0)\rangle$  as follows  $|\psi(t)\rangle = \hat{U}(t) |\psi(0)\rangle$
- For the walk on the line when  $k = 2$ , we get exactly the same distribution as the discrete case (figure 2), despite their very different constructions

## PROBABILITY DISTRIBUTIONS

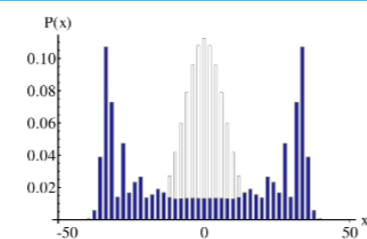


Figure 2: Probability distributions of the discrete and continuous quantum and classical walks on the line, for  $t = 50$

## GROVER'S ALGORITHM

What Is It?

- Originally discovered by Lov Grover in 1996-1997 [5]
- It is a very simple, but powerful database search algorithm that makes use of the linear superposition of states
- Imagine that you have a list of  $N$  names, that are unsorted, and you only want one of those names. Grover's algorithm will find the given name “marked state”, in a computational time  $t = \mathcal{O}(\sqrt{N})$
- This is remarkable, as the equivalent classical algorithm takes a time of  $t = \mathcal{O}(N)$ ! This means that if you had a 1,000,000 unsorted names, Grover's algorithm, on average, would find the name after searching 1,000 names. Whereas the classical case, would on average, have to search through at least 500,000 names!

Set Up

- We begin by preparing all states (elements of our database) in a linear superposition  $|\psi\rangle = \sum_{i=1}^N \frac{1}{\sqrt{N}} |i\rangle$ , (figure 3a) where  $\mu = \frac{1}{N}$
- Then we continually apply Grover's unitary operator  $\hat{G}$  recursively to an initial state  $|\psi(0)\rangle = |\psi\rangle \otimes |-\rangle$  and then make a measurement on the final state  $|\psi(t)\rangle$  after a number of steps  $t$ . Where  $|-\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$
- Grover's operator consists of an oracle  $\hat{O}$ , which acts as a function that shifts the phase of the marked state by  $\pi$ , whilst leaving the rest unchanged (figure 3b) and a diffusion transform  $\hat{D}$ , that inverts our marked state about the mean (figure 3c)

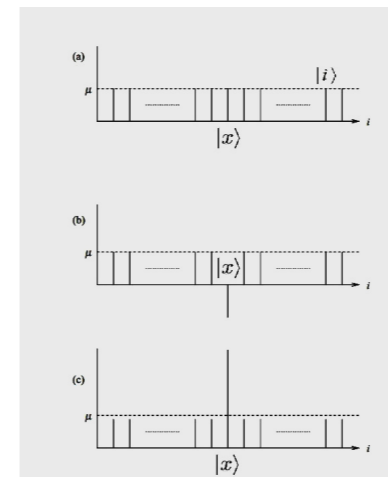


Figure 3: Amplitude amplification and Grover's algorithm at work

A Modified Version of Grover's Algorithm

- Due to the impressive properties of the algorithm, people had tried to apply Grover's algorithm directly to a spatial search, but unfortunately the algorithm reverts back to a computation time of  $\mathcal{O}(N)$  [6]. As in a spatial search there is an additional time cost as we move between memory locations
- Ambainis *et al.* [7] showed that we can use the construction of the discrete quantum walk to give us Grover's algorithm. In doing this they were able to get a computation time of  $\mathcal{O}(\sqrt{N} \log N)$  for 2-dimensions and  $\mathcal{O}(\sqrt{N})$  for 3 or more dimensions
- In using the walk for this simple case, we find that it not only enhanced computation time, but it also offered a new approach for constructing quantum algorithms

## EXPERIMENTAL IMPLEMENTATION

Quantum Computers

- To effectively implement the walks, we must be able to design a quantum computer
- There are several different types of experimental set ups, from nuclear magnetic resonance (NMR), optical photon computer, harmonic oscillator computer, ion traps and many others [8]
- We specifically focus on ion trapping, due to its similarities to the discrete walk and recent successes [9]

Trapped Ion Experiment

- We begin with a single Beryllium ion  $\text{Be}^+$  confined in a coaxial resonator radio frequency ion trap
- Next we create the unitary equation (1), by applying a sequence of four Raman beam pulses, to create a superposition state  $|\psi\rangle = \frac{1}{\sqrt{2}}(|\alpha\rangle | \downarrow \rangle + |-\alpha\rangle | \uparrow \rangle)$ , where  $|\alpha\rangle$  are the coherent states
- Then we measure the internal state of the ion, with the measurement operator  $\hat{M} = e^{\pm i\hat{p}\sigma_2}$ , where  $i\hat{p}\sigma_2$  is the Hamiltonian, which Hamiltonian we use is dependent on the internal state of the system. We then measure the internal state again
- If decoherence has effected the ion, we revert back to the classical walk. If there is no decoherence, then we get the quantum walk.

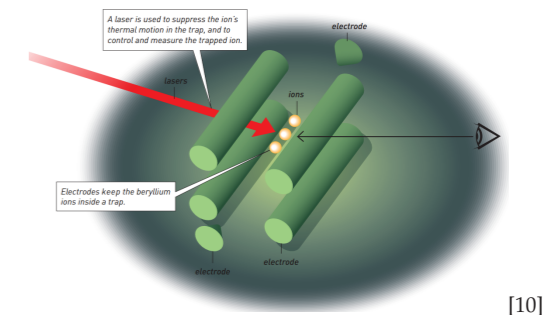


Figure 4: Ion trap. Using electromagnetic fields and light to confine, control, and measure the quantum state of beryllium ions.

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